Quantum Simulation of Gauge Theory

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with Henry Lamm and Yukari Yamauchi (NuQS Collaboration) Based on 1806.06649 and 190x.xxxxx

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Overview

Overview of quantum algorithms





Simulations of field theories

Simulating D_4 gauge theory



Quantum Chromodynamics



Lattice **QCD**



$$Z = \int_{SU(3)} \mathrm{d}U \ e^{-\int L}$$

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But we can't do...

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- Finite fermion density (sign problem)
- Viscosity (real-time evolution)

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Quantum computers promise all this!

From Bit to Qubit

|0
angle , |1
angle ...

From Bit to Qubit



This is a spin from QM.



A Quantum Computer

Physically, there's a Hilbert space:

 $\mathcal{H}=\mathcal{H}_1\otimes\mathcal{H}_1\otimes\cdots$

When we measure, we collapse into one of the 2^N states in the "fiducial basis".

 $|\Psi\rangle \rightarrow |0101010\rangle$

Gates

Arbitrary one-qubit gates are 'easy' – can be constructed from Hadamard and $\frac{\pi}{8}$ -gate.

$$H=rac{1}{\sqrt{2}} egin{pmatrix} 1&1\ 1&-1 \end{pmatrix}$$
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Controlled-not (CNOT) is a 2-qubit gate.



 $egin{aligned} |00
angle &\mapsto |00
angle \ |01
angle &\mapsto |01
angle \ |10
angle &\mapsto |11
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 $T \equiv \mathcal{R}_{\pi/4}$

What is a classical computer?

A quantum computer constantly being measured in the fiducial basis.

$$|01\rangle\,$$
 is okay — $[|00\rangle+|11\rangle]\,$ is not

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$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

Classical Algorithms Are Quantum Algorithms

Any classical circuit can be made into a quantum circuit!

Example: two-bit adder

 $\begin{array}{l} \left| 00 \right\rangle \left| 00 \right\rangle \rightarrow \left| 00 \right\rangle \left| 00 \right\rangle \\ \left| 01 \right\rangle \left| 00 \right\rangle \rightarrow \left| 01 \right\rangle \left| 01 \right\rangle \\ \left| 10 \right\rangle \left| 00 \right\rangle \rightarrow \left| 10 \right\rangle \left| 01 \right\rangle \\ \left| 11 \right\rangle \left| 00 \right\rangle \rightarrow \left| 11 \right\rangle \left| 10 \right\rangle \end{array}$



In general, given a classical function f(x), we can implement:

 $\ket{x}\ket{0} \rightarrow \ket{x}\ket{f(x)}$

Inverting Circuits

$$(AB)^{-1} = B^{-1}A^{-1}$$
$$-H + T^{\dagger}$$
$$-T + H + H$$

As long as we have the inverse of each individual gate...

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Easy version of real-time evolution: $U(t) = e^{-iHt}$, with H diagonal.

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{-i\theta} \end{pmatrix}$$
$$|11\rangle \rightarrow e^{-i\theta} |11\rangle$$
$$|0\rangle \longrightarrow \mathcal{R}_{\theta} \longrightarrow$$

Goal: e^{-iHt}

$$e^{-i(H_1+H_2)\epsilon} \approx e^{-iH_1\epsilon}e^{-iH_2\epsilon}$$

- 1. Split H into tiny pieces
- 2. Diagonalize each piece
- 3. Hit with phase gates

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- 3. Hit with phase gates

On a classical computer: $2^N \times 2^N$ matrices.

On a quantum computer: N qubits and $\propto t/\epsilon$ gates.

Example: Coupled Spins

 $H = \overbrace{\sigma_z(1)\sigma_z(2)}^{H_V} + \overbrace{\mu(\sigma_x(1) + \sigma_x(2))}^{H_K}$

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$$e^{-iHt} \approx \left(e^{-iH_V\epsilon}e^{-iH_K\epsilon}\right)^{t/\epsilon}$$

 H_V is diagonal!

$$egin{pmatrix} 1 & 0 & 0 & 0 \ 0 & e^{i\epsilon} & 0 & 0 \ 0 & 0 & e^{i\epsilon} & 0 \ 0 & 0 & 0 & 1 \end{pmatrix}$$

 H_K is diagonalized by the Hadamard gate.

$$-H - \mathcal{R}_{\epsilon} - H -$$
$$-H - \mathcal{R}_{\epsilon} - H -$$

Real-Time Evolution for Nonlinear Response



$$|\Psi\rangle \rightarrow e^{-iH_0T} |\Psi\rangle \rightarrow e^{-i(H_0+H_B)t} e^{-iH_0T} |\Psi\rangle \rightarrow \mathsf{Measure!}$$

Real-Time Evolution for Nonlinear Response



$$|\Psi\rangle \rightarrow e^{-iH_0T} |\Psi\rangle \rightarrow e^{-i(H_0+H_B)t} e^{-iH_0T} |\Psi\rangle \rightarrow \text{Measure}!$$

This gives $\langle \mathcal{O}(t) \rangle$. What about $\langle \mathcal{O}(t) \mathcal{O}(0) \rangle$?

Want linear response? Take a derivative!

Real-Time Evolution for Linear Response

Want linear response? Take a derivative!

$$H(t) = H_0 + \epsilon \delta(t) H'$$

- 1. Evolve briefly with $H' = \mathcal{O}$. $e^{-i\mathcal{O}\epsilon} |\Psi\rangle$
- 2. Perform normal time-evolution. $e^{-iH_0t}e^{-i\mathcal{O}\epsilon}\left|\Psi\right\rangle$
- 3. Measure \mathcal{O} .

$$\frac{\partial}{\partial \epsilon} \left\langle e^{iHt} \mathcal{O} e^{-iHt} \right\rangle = \operatorname{Im} \left\langle \mathcal{O}(t) \mathcal{O}(0) \right\rangle$$

Alternatives include: Roggero and Carlson 1804.01505; Pedernales et al. 1401.2430.

Parton Distribution Functions



$$f_q(\xi) = \int_{\infty}^{\infty} \frac{\mathrm{d}t}{2\pi} e^{-it\xi(n\cdot P)} \langle P | \bar{\psi}_q(tn^{\mu}) \frac{\hbar}{2} W_n \psi_q(0) | P \rangle$$

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Naive method: couple to a large thermal bath.

$$H = H_0 + H_{\rm bath} + H_{\rm int}$$

If $H_{\rm bath}$ is well-understood (easily arranged), we can prepare it cold, and then time-evolve.

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More sophisticated: Spectral Combing (1709.08250), Quantum Adiabatic Algorithm (quant-ph/0001106) both require e^{-iHt} .

Hybrid classical/quantum methods don't: 1806.06649.

The Hamiltonian of a free particle moving on G = SU(3):

$$H = -\nabla^2$$

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The Hilbert space is $\mathbb{C}G$: one basis state for every $U \in G$.

$$\left| \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \right\rangle, \left| \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & i \\ -i & 0 & 0 \end{pmatrix} \right\rangle, \left| \begin{pmatrix} 0 & 0 & 1 \\ \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} & 0 \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} & 0 \end{pmatrix} \right\rangle$$

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The Hamiltonian is diagonal in Fourier space. (For R, momentum space.)

Hamiltonian Lattice Gauge Theory



$$H = \frac{1}{g^2} \left[\sum_{L} \nabla_{L}^2 + \sum_{P} \operatorname{Re} \operatorname{Tr} P \right]$$

Gauge Symmetry



$$U_{ij}\mapsto V_j U_{ij} V_i^{\dagger}$$

$${
m Tr} \; U_{14}^{\dagger} U_{45}^{\dagger} U_{25} U_{12} \
ightarrow {
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m Tr} \; U_{14}^{\dagger} U_{45}^{\dagger} U_{45} U_{45} V_{45} V_{45$$

The Hamiltonian is gauge-invariant

$$H = \frac{1}{g^2} \left[\sum_{\langle ij \rangle} \nabla_{ij}^2 + \sum_P \operatorname{Re} \operatorname{Tr} P \right]$$

The Hilbert Space





The Hilbert Space



Only Gauge-Invariant States Allowed!

$$|U_{12}\rangle \qquad \int \mathrm{d}V_1 \mathrm{d}V_2 \left|V_2^{\dagger}U_{12}V_1\right\rangle$$
The Hilbert Space



$$\mathcal{H} = \mathbb{C} G \otimes \mathbb{C} G \otimes \cdots$$

Only Gauge-Invariant States Allowed!

$$\left| \begin{array}{c} \int \mathrm{d} V_1 \mathrm{d} V_2 \left| V_2^{\dagger} U_{12} V_1 \right\rangle \right.$$

The Hilbert Space



Only Gauge-Invariant States Allowed!

$$\left| \begin{array}{c} \int \mathrm{d}V_1 \mathrm{d}V_2 \right| V_2^{\dagger} U_{12} V_1 \right\rangle$$

Here's a projection operator:

$$P | U_{12} \cdots \rangle = \int (\mathrm{d} V_1 \mathrm{d} V_2 \cdots) | V_2^{\dagger} U_{12} V_1 \cdots \rangle$$

Trotterization

 $H = \frac{1}{g^2} \left[\sum_{L}^{H_{\mathcal{K}}} \nabla_L^2 + \sum_{P}^{H_{\mathcal{V}}} \operatorname{Re} \operatorname{Tr} P \right]$

Trotterization

$$H = \frac{1}{g^2} \left[\underbrace{\sum_{L}^{H_K} \nabla_L^2}_{L} + \underbrace{\sum_{P}^{H_V} \operatorname{Re} \operatorname{Tr} P}_{P} \right]$$

Kinetic

Potential

One link only Diagonal in Fourier space Mutually commuting terms Four links Diagonal (in our basis) Mutually commuting terms

$$e^{-iHt} \approx \left[\left(e^{-i\nabla_1^2 \epsilon} e^{-i\nabla_2^2 \epsilon} \cdots \right) \left(e^{-i\epsilon \operatorname{Re}\operatorname{Tr} P_1} e^{-i\epsilon \operatorname{Re}\operatorname{Tr} P_2} \cdots \right) \right]^{t/\epsilon}$$

Hilbert Space on a Quantum Computer

Classical algorithms \longrightarrow quantum algorithms!

 $SU(3) \longleftrightarrow \{0000, 0001, 0010, \cdots\}$

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That's a basis for $\tilde{\mathcal{H}}_1$! An "SU(3)-register" — the analog of a classical variable.

Now that we have $\mathcal{H}_1\sim\tilde{\mathcal{H}}_1,$ we can construct the full $\mathcal H$ on a quantum computer.

$$|U_{12}\rangle |U_{23}\rangle \cdots \in \tilde{\mathcal{H}}_1 \otimes \tilde{\mathcal{H}}_1 \otimes \cdots$$

The set of physical states is a linear subspace.

$$H = \frac{1}{g^2} \left[\sum_{L} \nabla_L^2 + \underbrace{\sum_{P} \operatorname{Re} \operatorname{Tr} P}_{P} \right]$$

We need an operator:

$$\mathcal{U}(heta)\ket{A}\ket{B}\ket{C}\ket{D} = e^{-i heta\operatorname{\mathsf{Re}}\operatorname{\mathsf{Tr}}(ABCD)}\ket{A}\ket{B}\ket{C}\ket{D}$$

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$$\begin{split} \mathcal{U}(\theta) \ket{A} \ket{B} \ket{C} \ket{D} &= e^{-i\theta \operatorname{\mathsf{Re}} \operatorname{\mathsf{Tr}}(ABCD)} \ket{A} \ket{B} \ket{C} \ket{D} \\ & \ket{A} \ket{B} \ket{C} \ket{D} \end{split}$$

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 $\mathcal{U}(\theta) |A\rangle |B\rangle |C\rangle |D\rangle = e^{-i\theta \operatorname{Re} \operatorname{Tr}(ABCD)} |A\rangle |B\rangle |C\rangle |D\rangle$ $|A\rangle |B\rangle |C\rangle |D\rangle$ $\rightarrow |A\rangle |B\rangle |C\rangle |D\rangle$ $\rightarrow \cdots \rightarrow |A\rangle |B\rangle |C\rangle |ABCD\rangle$

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$$H = \frac{1}{g^2} \left[\sum_{L}^{H_K} \nabla_L^2 + \sum_{P} \operatorname{Re} \operatorname{Tr} P \right]$$

Diagonal in "momentum basis". Need to perform a Quantum Fourier Transform. This is a Fourier transform of $\Psi(x)$:

$$|\Psi
angle = \sum_{x} \Psi(x) |x
angle$$

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For \mathbb{Z}_2 : σ_x is diagonal in Fourier space, and H performs QFT. For \mathbb{R} :

$${\sf F} \ket{x} = \sum_p e^{-i x p} \ket{p}$$

For *SO*(3): QFT decomposes into spherical harmonics So: QFT, then phase gate (on a single link!), then QFT.

The Dihedral Group D₄



As a matrix group, $D_4 < U(2)$.

$$\left\langle \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} i & 0 \\ 0 & -i \end{pmatrix} \right\rangle$$

The D₄-Register

 $|D_4| = 8$, so we want 3 qubits.





Inversion

$\mathsf{Classical\ circuits} \longrightarrow \mathsf{quantum\ circuits!}$

$$\begin{array}{ll} |000\rangle \rightarrow |000\rangle & |100\rangle \rightarrow |100\rangle \\ |001\rangle \rightarrow |011\rangle & |101\rangle \rightarrow |101\rangle \\ |010\rangle \rightarrow |010\rangle & |110\rangle \rightarrow |110\rangle \\ |011\rangle \rightarrow |001\rangle & |111\rangle \rightarrow |111\rangle \end{array}$$

$$X=egin{pmatrix} 0&1\1&0 \end{pmatrix} \leftrightarrow \ket{1} \end{pmatrix}$$



Multiplication

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 $\mathcal{U}_{ imes} \ket{U} \ket{V} = \ket{U} \ket{UV}$

Because G is a group, this operation is unitary.

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Because G is a group, this operation is unitary.



Multiplication and inversion let us construct a plaquette:

$$\mathcal{U}_{P} \ket{U_{12} \cdots} \ket{\mathbf{1}} = \ket{U_{12} \cdots} \ket{P}$$

Trace

$$\mathcal{U}_{\mathsf{Tr}}(heta) \ket{U} = e^{i heta \operatorname{\mathsf{Re}} \operatorname{\mathsf{Tr}} U} \ket{U}$$

Only two elements of D_4 have non-zero trace.

$$\mathsf{Tr} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = 2 \qquad \qquad \mathsf{Tr} \begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix} = -2$$

These correspond to the states $|000\rangle$ and $|010\rangle$.



Fourier Transform

$$\begin{pmatrix} \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & 0.0 & -\frac{1}{2} & 0.0 & \frac{1}{2} & 0.0 & -\frac{1}{2} & 0.0 \\ 0.0 & -\frac{1}{2} & 0.0 & -\frac{1}{2} & 0.0 & \frac{1}{2} & 0.0 & -\frac{1}{2} \\ 0.0 & \frac{1}{2} & 0.0 & -\frac{1}{2} & 0.0 & \frac{1}{2} & 0.0 \end{pmatrix} \end{pmatrix}$$

Fourier Transform

$$\begin{pmatrix} \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} \\ \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} & -\frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}} \\ \frac{1}{2} & 0.0 & -\frac{1}{2} & 0.0 & \frac{1}{2} & 0.0 & -\frac{1}{2} & 0.0 \\ 0.0 & -\frac{1}{2} & 0.0 & \frac{1}{2} & 0.0 & \frac{1}{2} & 0.0 & -\frac{1}{2} \\ 0.0 & \frac{1}{2} & 0.0 & -\frac{1}{2} & 0.0 & \frac{1}{2} & 0.0 & \frac{1}{2} & 0.0 \\ \frac{1}{2} & 0.0 & -\frac{1}{2} & 0.0 & -\frac{1}{2} & 0.0 & \frac{1}{2} & 0.0 \end{pmatrix}$$

Yikes! This is a job for a computer...



The **global** Hilbert space is large, and can't be treated classically. The **local** Hilbert space is small, so **it can be treated classically**.

The diagonalized kinetic operator is



 $\mathcal{H} = \mathbb{C} D_4 \otimes \mathbb{C} D_4 \otimes \cdots \text{ maps onto } 3 \times L \text{ qubits}$

- 1. Prepare initial state (somehow)
- 2. Trace circuit on all plaquettes
- 3. QFT on all links
- 4. Phase gate on the most-significant qubit of all links
- 5. QFT^{\dagger} on all links
- 6. Repeat 2-5 to get the desired t
- 7. Measure; look at something gauge-invariant

Requires \sim 300 gates per time step.

Real-Time Evolution



A Comment on Trotterization

Deriving time evolution

$$U = e^{-iHt}$$
$$= \left(e^{-i(H_1 + H_2)\delta}\right)^{t/\delta}$$
$$\approx \left(e^{-iH_1\delta}e^{-iH_2\delta}\right)^{t/\delta}$$

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Deriving time evolution

l

$$egin{aligned} \mathcal{J} &= e^{-iHt} \ &= \left(e^{-i(H_1+H_2)\delta}
ight)^{t/\delta} \ &pprox \left(e^{-iH_1\delta}e^{-iH_2\delta}
ight)^{t/\delta} \end{aligned}$$

Deriving the euclidean lattice

$$Z = \operatorname{Tr} e^{-\beta H}$$
$$= \operatorname{Tr} \left(e^{-\delta(H_1 + H_2)} \right)^{\beta/\delta}$$
$$\approx \operatorname{Tr} \left(e^{-\delta H_1} e^{-\delta H_2} \right)^{\beta/\delta}$$

It's the same procedure!

A Comment on Trotterization

Deriving time evolution

$$U = e^{-iHt} \qquad Z = \operatorname{Tr} e^{-\beta H}$$
$$= \left(e^{-i(H_1 + H_2)\delta}\right)^{t/\delta} \qquad = \operatorname{Tr} \left(e^{-\delta(H_1 + H_2)}\right)^{\beta/\delta}$$
$$\approx \left(e^{-iH_1\delta}e^{-iH_2\delta}\right)^{t/\delta} \qquad \approx \operatorname{Tr} \left(e^{-\delta H_1}e^{-\delta H_2}\right)^{\beta/\delta}$$

It's the same procedure!

Deriving the euclidean lattice

Physics on the classical lattice should be similar to physics on the quantum lattice.

Coupling Constants on the Cheap!

$$H = \frac{1}{g^2} \left[a^3 \sum_{\langle ij \rangle} \nabla_{ij}^2 + a^2 \sum_P \operatorname{Re} \operatorname{Tr} P \right] + a^3 m \sum_i \bar{\psi}_i \psi_i + a^2 \sum_{\langle ij \rangle} \bar{\psi}_i \psi_j$$

We want $m_{\pi} \approx 135 \text{MeV}$. What should *a*, *m*, and *g* be?

Coupling Constants on the Cheap!

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We want $m_{\pi} \approx 135 \text{MeV}$. What should *a*, *m*, and *g* be?

Use the euclidean lattice!

Spatial discretization and finite volume effects are the same: desired couplings can be found classically.

The Present

The Near Future





Current best ~ 10 qubits 10 gates

"Quantum supremacy" 50 qubits 40 gates

Needed (for 10³ lattice)

 $\overbrace{10^3\times3}^{links}\times \overbrace{16\times18}^{\textit{SU}(3)\text{-register}}\sim 10^6 \text{ qubits}$

Why bother?

Concrete staring point for more efficient algorithms Are large-scale quantum processors worth building?
Why bother?

Concrete staring point for more efficient algorithms Are large-scale quantum processors worth building?

Inspiration for new classical algorithms

A quantum-inspired classical algorithm for recommendation systems

Ewin Tang

July 13, 2018

Abstract

A recommendation system suggests products to users based on data about user preferences. It is typically modeled by a problem of completing an $m \times n$ matrix of small rank k. We give the first classical algorithm to produce a recommendation in $O(\text{poly}(k) \operatorname{poly}(g(m, n))$ time, which is an exponential improvement on previous algorithms that run in time linear in m and n. Our strategy is inspired by a quantum algorithm by Kerenidis and Prakash: like the quantum algorithm, instead of recom-

- How few qubits can we get away with?
- Can error correction/tolerance be done more cheaply?
- How to prepare QCD ground state?
- What algorithms *don't* involve e^{-iHt} ?

Outline

Simulating QCD

Overview of Quantum Computing

Straightforward Algorithms

A Crucial Building Block

Simulating Gauge Theories

Bits and Pieces — D₄

Simulated Simulations

Choosing Coupling

Future Work